

# Simply Typed Lambda Calculus (STLC)

$\tau ::= \text{unit} \mid \tau \rightarrow \tau \mid \tau \times \tau \mid \tau + \tau \mid \text{void}$

$e ::= x \mid () \mid e \ e \mid \lambda x: \tau. e \mid (e, e) \mid \text{fst } e \mid \text{snd } e \mid \text{inl } e \mid \text{inr } e$   
 case  $e$  of  $\{x.e; y.e\} \mid \text{abort } e$   
 $\Gamma \vdash e: \tau$

$$\frac{\Gamma(x) = \tau \quad (T1) \quad \Gamma \vdash () : \text{unit} \quad (T2)}{\Gamma \vdash x: \tau} \quad \frac{\Gamma \vdash e_1: \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash e_2: \tau_1 \quad (T-3)}{\Gamma \vdash e_1 \ e_2: \tau_2}$$

$$\frac{\Gamma, x: \tau \vdash e: \tau' \quad (T4)}{\Gamma \vdash \lambda x: \tau. e: \tau \rightarrow \tau'}$$

$$\frac{\Gamma \vdash e_1: \tau_1 \quad \Gamma \vdash e_2: \tau_2 \quad (T-5)}{\Gamma \vdash (e_1, e_2): \tau_1 \times \tau_2} \quad \frac{\Gamma \vdash e: \tau_1 \times \tau_2 \quad (T-6)}{\Gamma \vdash \text{fst } e: \tau_1} \quad \frac{\Gamma \vdash e: \tau_1 \times \tau_2 \quad (T-7)}{\Gamma \vdash \text{snd } e: \tau_2}$$

$\text{dims} \equiv \text{int} \times \text{int}$   
 $\quad \quad \quad \uparrow \quad \quad \uparrow$   
 $\quad \quad \quad x \quad \quad y$

$\text{person} \equiv \text{string} \times \text{int}$   
 $\quad \quad \quad \uparrow \quad \quad \uparrow$   
 $\quad \quad \quad \text{name} \quad \text{age}$

$\text{test} \equiv \text{negative} \mid \text{positive} \leftarrow \text{enum (e.g. Java)}$

$\text{lolla}: \text{proof} \rightarrow \text{bool}$

$\text{proof} \equiv \text{vaccine of string} \times \text{int}$   
 or  $\text{test of test}$

$\text{proof} \equiv (\text{string} \times \text{int}) + \text{test}$

$\text{test} \equiv \text{unit} + \text{unit}$

$\text{neg\_test} \equiv \text{inl } \text{negative}$

$\text{negative} \equiv \text{inl } ()$

$\text{pfizer\_vax} \equiv \text{inl } ("Pfizer", 2)$

$\text{positive} \equiv \text{inr } ()$

$\text{lolla} \equiv \lambda p: \text{proof}. \text{case } p \text{ of}$

$\{ \text{vax}. \text{snd } \text{vax} = 2 \parallel \text{fst } \text{vax} = "J \ \& \ J";$   
 $\text{test}. \text{test} = \text{negative} \}$

$$\frac{\Gamma \vdash e : \tau_1}{\Gamma \vdash \text{inl } e : \tau_1 + \tau_2} \quad (\text{T-8}) \quad \frac{\Gamma \vdash e : \tau_2}{\Gamma \vdash \text{inr } e : \tau_1 + \tau_2} \quad (\text{T-9})$$

$$\frac{\Gamma \vdash e : \tau_1 + \tau_2 \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case } e \text{ of } \{x. e_1, y. e_2\} : \tau} \quad (\text{T-10})$$

Notice: Binary product  $\tau_1 \times \tau_2$   
 1 way to create (pair)  
 2 things to do with it (fst, snd)

n-ary product  $\tau_1 \times \dots \times \tau_n$   
 1 way to create (tuple)  
 n things to do (projections)

Binary sum  $\tau_1 + \tau_2$   
 2 ways to create (inl, inr)  
 1 thing to do (case)

n-ary sum  $\tau_1 + \dots + \tau_n$   
 n ways to create  
 1 thing to do

0-ary (nullary) product  
 1 way to create (supply nothing) = Unit!  
 0 things to do

nullary sum: 0 ways to create, 1 thing to do "void"

$$\frac{\Gamma \vdash e : \text{void}}{\Gamma \vdash \text{abort } e : \tau} \quad (\text{T-11})$$

$$\frac{v \text{ val}}{(\lambda x:\tau) v \mapsto [v/x]c}$$

$$(\lambda x.c) \text{ val}$$

$$\frac{() \text{ val} \quad \frac{v_1 \text{ val} \quad v_2 \text{ val}}{(v_1, v_2) \text{ val}}}{() \text{ val}}$$

$$\frac{e_1 \mapsto e_1'}{(e_1, e_2) \mapsto (e_1', e_2)}$$

$$\frac{v_1 \text{ val} \quad e_2 \mapsto e_2'}{(v_1, e_2) \mapsto (v_1, e_2')}$$

$$\frac{\dots e \mapsto e'}{fst \cdot e \mapsto fst \cdot e'}$$

$$\frac{v_1 \text{ val} \quad v_2 \text{ val}}{fst (v_1, v_2) \mapsto v_1}$$

$$\frac{e \mapsto e'}{inl e \mapsto inl e'}$$

$$\frac{v \text{ val}}{inl v \text{ val}}$$

$e \mapsto e'$

case  $e$  of  $\{x.e_1; y.e_2\} \mapsto$  case  $e'$  of  $\{x.e_1; y.e_2\}$

case  $inl v$  of  $\{x.e_1; y.e_2\} \mapsto [v/x].e_1$

case  $inr v$  of  $\{x.e_1; y.e_2\} \mapsto [v/y].e_2$

Progress: If  $\Gamma \vdash e:\tau$  then  $e \text{ val}$  or  $e \mapsto e'$

Preservation: If  $\Gamma \vdash e:\tau$  and  $e \mapsto e'$  then  $\Gamma \vdash e':\tau$ .

Fun fact: If  $\vdash e:\tau$  then  $\exists v \text{ st. } v \text{ val and } e \mapsto^* v$   
(remember: not true for untyped  $\lambda$ -calculus)

$$\frac{\Gamma, x:\tau \vdash x:\tau \quad \Gamma, x:\tau \vdash x x:\tau}{\Gamma, x:\tau \vdash x x:\tau}$$

$\tau = \tau \rightarrow \tau$  WAT?

$$\Gamma \vdash (\lambda x. x x) (\lambda x. x x):$$