Lecture 3

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Last time: Type system to reject programs like
`1 "Hello" ^ "World" 1`

Type safety: "Well-typed programs can't go wrong"
-Robin Milner

2 components:
Progress: If e is well-typed, it's a value or can take a step.

Preservation: If a well-typed exp. takes a step, still well-typed w/ the same type.

$e_1 :\tau \mapsto e_2 :\tau \mapsto e_3 :\tau \mapsto \ldots \mapsto v$

Preservation: If $e:\tau$ and $e \mapsto e'$ then $e':\tau$
Pf: By induction on the derivation of $e \mapsto e'$

S-1 By inversion on T-3, $\tau = int$.
By T-1, $\overline{n_1 + n_2} : int$.
S-2 By inversion on T-4, $\tau = string$. By T-2, $"s_1 s_2" : string$.
S-3 By inversion on T-5, $\tau = int$. By T-1, $\overline{|s|} : int$.
S-4 By inversion on T-3, $\tau = int$, $e_1 : int$, $e_2 : int$
By IH, $e_1' : int$.
By T-3, $e_1' + e_2 : int$.
S-5. By inversion on T-3, $\tau = int$, $e_2 : int$. $\overline{n_1} : int$
By IH, $e_2' : int$.
By T-3, $\overline{n_1} + e_2' : int$.
S-6, S-7, S-8 similar to above. □

Lemma: Canonical Forms
1. If e val and $e:int$, then $e = \bar{n}$ for some n.
2. If e val and $e:string$, then $e = "s"$ for some s.

Pf: The only rules that can derive e val are V-1 and V-2.
If V-1, then $e = \bar{n}$ and $e:int$.
If V-2, then $e = "s"$ and $e:string$. □

Progress: If $e:\tau$, then e val or there exists e' s.t. $e \mapsto e'$.
Pf: By induction on the derivation of $e:\tau$.
T-1 $\bar{n}$ val by V-1
T-2 "s" val by V-2
T-3 Then $\tau = int$, $e = e_1 + e_2$, $e_1:int$, and $e_2:int$.
By IH, e_1 val or $e_1 \mapsto e_1'$.
- e_1 val. By CF, $e_1 = \bar{n}_1$ for some n_1.
  By IH, e_2 val or $e_2 \mapsto e_2'$.
  - e_2 val. By CF, $e_2 = \bar{n}_2$ for some n_2.
    By S+1, $\bar{n}_1 + \bar{n}_2 \mapsto \overline{n_1 + n_2}$.
  - $e_2 \mapsto e_2'$. By S-5, $\bar{n}_1 + e_2 \mapsto \bar{n}_1 + e_2'$.
- $e_1 \mapsto e_1'$. By S-4, $e_1 + e_2 \mapsto e_1' + e_2$.

T-4, T-5 similar to above. □