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CWID

## Exam <br> 2

Dec 7th, 2018

## CS525 - Final Exam Solutions

Please leave this empty! 1 $\square$
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$\square$ $7 \square$ $\operatorname{Sum} \square$

## Instructions

- Things that you are not allowed to use
- Personal notes
- Textbook
- Printed lecture notes
- Phone
- Calculator
- The exam is $\mathbf{1 2 0}$ minutes long
- Multiple choice questions are graded in the following way: You get points for correct answers and points subtracted for wrong answers. You do not have to answer every subquestion. The minimum points for each questions is $\mathbf{0}$. For example, assume there is a multiple choice question with 6 answers - each may be correct or incorrect - and each answer gives 1 point. If you answer 3 questions correct and 3 incorrect you get 0 points. If you answer 4 questions correct and 2 incorrect you get 2 points. If you answer 3 questions correct, 1 incorrect, and do not answer 2 , then you get $3-1=2$ points. .
- For your convenience the number of points for each part and questions are shown in parenthesis.
- There are 7 parts in this exam (100 points total)

1. SQL (26)
2. Relational Algebra (15)
3. Index Structures (20)
4. I/O Estimation (12)
5. Result Size Estimation (12)
6. Schedules (15)

## Part 1 SQL (Total: 26 Points)

Consider the following tvshow database schema and instance:
show

| title | genre | rating | costPerEpisode |
| :---: | :---: | :---: | :---: |
| Game of Thrones | Fantasy | 7.8 | $32,000,000$ |
| Monk | Crime | 6.2 | $1,500,000$ |
| Dick Tracy | Crime | 4.3 | 800,000 |
| Ancient Aliens | Conspiracy | 2.2 | 200,000 |

## channel

| cName | category | networth |
| :---: | :---: | :---: |
| Fox News | news | $4,000,000,000$ |
| HBS | subscription | $3,000,000,000$ |
| History | documentary \& aliens | $400,000,000$ |
| MSMBC | news | $5,000,000,000$ |

## schedule

| date | tId | cName | show |
| :---: | :---: | :---: | :---: |
| 2013-01-01 | 11 | HBS | Monk |
| $2013-01-01$ | 11 | History | Ancient Aliens |
| $2013-01-01$ | 12 | Fox News | Dick Tracy |
| $2013-01-02$ | 1 | HBS | Game of Thrones |

## timeslot

| tId | start | end | avgViewers |
| :---: | :---: | :---: | :---: |
| 1 | $00: 00$ | $02: 00$ | 500,000 |
| 2 | $02: 00$ | $04: 00$ | 300,000 |
| 3 | $04: 00$ | $06: 00$ | 50,000 |
| 4 | $06: 00$ | $08: 00$ | $1,000,000$ |
| 5 | $08: 00$ | $10: 00$ | $3,000,000$ |
| 6 | $10: 00$ | $12: 00$ | $2,000,000$ |
| 7 | $12: 00$ | $10: 00$ | $1,000,000$ |
| 8 | $14: 00$ | $16: 00$ | $1,500,000$ |
| 9 | $16: 00$ | $18: 00$ | $4,000,000$ |
| 10 | $18: 00$ | $20: 00$ | $6,000,000$ |
| 11 | $20: 00$ | $22: 00$ | $4,500,000$ |
| 12 | $22: 00$ | $00: 00$ | $2,300,000$ |

## episodes

| showtitle | eTitle | season | episode |
| :---: | :---: | :---: | :---: |
| Monk | Monkish | 1 | 1 |
| Monk | Hello World | 1 | 2 |
| Dick Tracy | The Murder Popoerder | 15 | 6 |

## Hints:

- When writing queries do only take the schema into account and not the example data given here. That is your queries should return correct results for all potential instances of this schema.
- Attributes with black background form the primary key of an relation. For example, title is the primary key of relation show.
- The attributes tId of relation schedule is a foreign key to relation timeslot.
- The attributes $\mathbf{c N a m e}$ of relation schedule is a foreign key to relation channel.
- The attributes show of relation schedule is a foreign key to relation show.
- Attribute showtitle of relation episode is a foreign key to relation show.


## Question $1.1 \quad$ (4 Points)

Write an SQL statement that returns the names (cName) of the channels with a networth that is higher than the average networth of all channel.

## Solution

```
SELECT cName
FROM channel
WHERE networth > (SELECT avg(networth) FROM channel);
```


## Question $1.2 \quad(5$ Points)

Write an SQL query that computes the total cost of each TV show. The total cost is the costPerEpisode multiplied with the number of episodes (table episodes) for this show.

## Solution

```
SELECT title, costPerEpisode * count(*) AS totalCost
FROM show s, episodes e
WHERE s.title = e.showtitle
GROUP BY title, costPerEpisode;
```

Also ok to use join of course.

## Question 1.3 (5 Points)

Write an SQL query that returns for each show the cost it has per viewer (the total number of viewers per show is the sum of avgViewers over all timeslots where the show is shown).

## Solution

```
SELECT title, costPerEpisode / sum(avgViewers) AS costPerViewer
FROM show s, schedule d, timeslot t
WHERE s.title = d.show AND d.tId = t.tId
GROUP BY s.title
```


## Question 1.4 (6 Points)

Write an SQL query that returns the show with the highest amount of total viewers per category. For that calculate the number of viewers by summing up the average viewers for each time slot when the show is shown (table schedule).

## Solution

```
WITH viewersPerChannel AS (
    SELECT s.cName, category, sum(avgViewers) AS totalViewers
    FROM channel c, schedule s, timeslot t
    WHERE c.cName = s.cName AND s.tId = t.tId
    GROUP BY s.cName, category
    )
SELECT cName, category
FROM viewersPerChannel v1
WHERE v1.totalViewers = (SELECT max(v2.totalViewers)
    FROM viewersPerChannel v2
    WHERE v2.category = v1.category);
```


## Question 1.5 (6 Points)

Write an SQL query that returns shows (their title and genre) with high costs (above $\$ 1,000,000$ per episode) that have low ratings ( 3.0 or lower), have less than $2,000,000$ viewers in total, and have run for more than 3 seasons.

## Solution

```
SELECT title, genre
FROM show s
WHERE rating < 3.0
    AND costPerEpisode > 1000000
    AND 2000000 < (SELECT sum(avgViewers)
                FROM timeslot t, schedule d
                WHERE t.tId = d.tId AND d.show = s.title)
    AND title IN (SELECT showtitle
                FROM episodes
                GROUP BY showtitle
                HAVING max(season) > 3)
```


## Part 2 Relational Algebra (Total: 15 Points)

## Question 2.1 Relational Algebra (4 Points)

Write a relational algebra expression over the schema from the SQL part that returns the titles of all episodes for a named "Monk". Use the bag semantics version of relational algebra.

## Solution

$$
\pi_{e T i t l e}\left(\sigma_{\text {showtitle }=\text { Monk }}(\text { episodes })\right)
$$

## Question 2.2 Relational Algebra (5 Points)

Write a relational algebra expression over the schema from the SQL part that returns the combined networth of all news channels. Use the bag semantics version of relational algebra.

## Solution

$$
\alpha_{\text {sum }(\text { networth })}\left(\sigma_{\text {category } y=\text { news }}(\text { channel })\right)
$$

## Question 2.3 Relational Algebra (6 Points)

Write a relational algebra expression over the schema from the SQL part that returns the average rating of shows per channel. Use the bag semantics version of relational algebra.

Solution
${ }_{c N a m e} \alpha_{\text {avg(rating) }}\left(\right.$ schedule $\bowtie_{\text {title=show }}$ show $\bowtie_{c N a m e=c N a m e}$ channel $)$

## Part 3 Index Structures (Total: 20 Points)

## Question 3.1 B+-tree Operations (20 Points)

Given is the $\mathrm{B}+$-tree shown below $(n=3)$. Execute the following operations and write down the resulting $B+$-tree after each step:
insert(55), insert(7), delete(13), delete(23), insert(3)

When splitting or merging nodes follow these conventions:

- Leaf Split: In case a leaf node needs to be split, the left node should get the extra key if the keys cannot be split evenly.
- Non-Leaf Split: In case a non-leaf node is split evenly, the "middle" value should be taken from the right node.
- Node Underflow: In case of a node underflow you should first try to redistribute and only if this fails merge. Both approaches should prefer the left sibling.



## Solution

insert(55)

insert(7)

delete(13)

delete(23)

insert(3)


## Part 4 I/O Cost Estimation (Total: 12 Points)

## Question 4.1 External Sorting (6 Points)

You have $M=101$ memory pages available and should sort a relation $R$ with $B(R)=2,000$ blocks. Estimate the number of I/Os necessary to sort $R$ using the external merge sort algorithm introduced in class.

## Solution

$$
\begin{aligned}
I O & =2 \cdot B(R) \cdot\left(1+\left\lceil\log _{M-1}\left(\frac{B(R)}{M}\right)\right\rceil\right) \\
& =2 \cdot 2,000 \cdot(1+1) \\
& =8,000
\end{aligned}
$$

## Question 4.2 External Sorting (6 Points)

You have $M=101$ memory pages available and should sort a relation $R$ with $B(R)=800,000$ blocks. Estimate the number of I/Os necessary to sort $R$ using the external merge sort algorithm introduced in class.

## Solution

$$
\begin{aligned}
I O & =2 \cdot B(R) \cdot\left(1+\left\lceil\log _{M-1}\left(\frac{B(R)}{M}\right)\right\rceil\right) \\
& =2 \cdot 800,000 \cdot(1+2) \\
& =4,800,000
\end{aligned}
$$

## Part 5 Result Size Estimations (Total: 12 Points)

Consider the table show relation from the SQL part with attributes title, genre, rating, and costPerEpisode. Attribute title is the primary key of this relation.

Given are the following statistics:

$$
\begin{aligned}
T(\text { show }) & =2,000 & & \\
V(\text { show }, \text { title }) & =2,000 & & \\
V(\text { show, genre }) & =10 & \min (\text { rating }) & =0.0
\end{aligned} \quad \max (\text { rating })=10.0
$$

## Question 5.1 Estimate Result Size (5 Points)

Estimate the number of result tuples for the query $q=\sigma_{\text {genre }=\text { Fantasy } \wedge \text { costPerEpisode } \geq 1,000,000 \text { (show) using the }}$ first assumption presented in class (values used in queries are uniformly distributed within the active domain).

## Solution

Calculate probability that a tuple fulfills the conditions using the independence assumption.

$$
\begin{gathered}
P(\text { genre }=\text { Fantasy } \wedge \text { costPerEpisode } \geq 1,000,000)=\frac{1}{10} \cdot \frac{2,099,999-1,000,000+1}{2,099,999-100,000+1}=\frac{1}{20}=0.05 \\
T(q)=T(\text { show }) * P(\text { genre }=\text { Fantasy } \wedge \text { costPerEpisode } \geq 1,000,000)=2,000 \cdot 0.05=10
\end{gathered}
$$

## Question 5.2 Estimate Result Size (7 Points)

Estimate the number of result tuples for the query $q=\sigma_{\text {genre }=C r i m e \vee g e n r e=N e w s ~}($ show $)$ using the first assumption presented in class.

## Solution

Since genre $=$ Crime and genre $=$ News are disjoint events we can sum up their probabilities to calculate $P($ genre $=$ Crime $\vee$ genre $=$ News $)$.

$$
\left.T(q)=\frac{2}{V(\text { show }, \text { genre })}\right) \cdot T(\text { show })=\frac{2}{10} \cdot 2,000=400
$$

We also accept the standard formula for disjunction even though it is not the right one to apply here.
$T(q)=\left(1-\left[\left(1-\frac{1}{V(\text { show, genre })}\right) \cdot\left(1-\frac{1}{V(\text { show, genre })}\right)\right]\right) \cdot T($ show $)=\left(1-\left[\left(1-\frac{1}{10}\right) \cdot\left(1-\frac{1}{10}\right)\right] \cdot 2,000=380\right.$

## Part 6 Schedules (Total: 15 Points)

## Question 6.1 Schedule Classes ( $15=5+5+5$ Points)

Indicate which of the following schedules belong to which class. Every correct answer is worth 1 point. Every incorrect answer results in 1 point being deducted. You are allowed to skip questions (0 points). Recall transaction operations are modelled as follows:

| $w_{1}(A)$ | transaction 1 wrote item $A$ |
| :---: | :--- |
| $r_{1}(A)$ | transaction 1 read item $A$ |
| $c_{1}$ | transaction 1 commits |
| $a_{1}$ | transaction 1 aborts |

$$
\begin{aligned}
& S_{1}=w_{1}(A), w_{4}(A), r_{2}(B), r_{3}(A), w_{1}(A), w_{4}(B), c_{1}, c_{2}, c_{3}, c_{4} \\
& S_{1}=r_{4}(B), w_{1}(B), w_{1}(D), c_{1}, r_{3}(D), w_{3}(C), c_{3}, r_{2}(C), w_{2}(A), r_{4}(A), c_{2}, c_{4} \\
& S_{3}=w_{4}(A), w_{1}(B), r_{1}(A), w_{3}(B), w_{4}(C), r_{2}(B), w_{2}(B), c_{2}, c_{1}, c_{3}, c_{4}
\end{aligned}
$$

| $S_{1}$ is recoverable | $\square$ no | $\square$ yes |
| :--- | :--- | :--- |
| $S_{1}$ is cascade-less | $\square$ no | $\square$ yes |
| $S_{1}$ is strict | $\square$ no | $\square$ yes |
| $S_{1}$ is conflict-serializable | $\square$ no |  |
| $S_{1}$ is 2 PL | $\square$ nos |  |
| yes |  |  |

$S_{3}$ is recoverable
$S_{3}$ is cascade-less
$S_{3}$ is strict
$S_{3}$ is conflict-serializable
$S_{3}$ is 2 PL
$S_{2}$ is recoverable
$S_{2}$ is cascade-less
$S_{2}$ is strict
$S_{2}$ is conflict-serializable
$S_{2}$ is 2 PL $\square$ no $\square$

