



CS425 – Summer 2016

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Chapter 4: Introduction to SQL

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Chapter 4: Introduction to SQL

- n Overview of the SQL Query Language
- n Data Definition
- n Basic Query Structure
- n Additional Basic Operations
- n Set Operations
- n Null Values
- n Aggregate Functions
- n Nested Subqueries
- n Modification of the Database



Textbook: Chapter 3



History

- n IBM Sequel language developed as part of System R project at the IBM San Jose Research Laboratory
- n Renamed Structured Query Language (SQL)
- n ANSI and ISO standard SQL:
 - | SQL-86, SQL-89, SQL-92
 - | SQL:1999, SQL:2003, SQL:2008
- n Commercial systems offer most, if not all, SQL-92 features, plus varying feature sets from later standards and special proprietary features.
 - | Not all examples here may work one-to-one on your particular system.



Data Definition Language

The SQL **data-definition language (DDL)** allows the specification of information about relations, including:

- n The schema for each relation.
- n The domain of values associated with each attribute.
- n Integrity constraints
- n And as we will see later, also other information such as
 - | The set of indices to be maintained for each relations.
 - | Security and authorization information for each relation.
 - | The physical storage structure of each relation on disk.



Domain Types in SQL

- n char(*n*).** Fixed length character string, with user-specified length *n*.
- n varchar(*n*).** Variable length character strings, with user-specified maximum length *n*.
- n int.** Integer (a finite subset of the integers that is machine-dependent).
- n smallint.** Small integer (a machine-dependent subset of the integer domain type).
- n numeric(*p,d*).** Fixed point number, with user-specified precision of *p* digits, with *n* digits to the right of decimal point.
- n real, double precision.** Floating point and double-precision floating point numbers, with machine-dependent precision.
- n float(*n*).** Floating point number, with user-specified precision of at least *n* digits.
- n** More are covered in Chapter 4.



Create Table Construct

- n An SQL relation is defined using the **create table** command:

```
create table  $r$  ( $A_1 D_1, A_2 D_2, \dots, A_n D_n,$   
                (integrity-constraint1),  
                ...,  
                (integrity-constraintk))
```

- | r is the name of the relation
 - | each A_i is an attribute name in the schema of relation r
 - | D_i is the data type of values in the domain of attribute A_i
- n Example:

```
create table instructor (  
    ID           char(5),  
    name        varchar(20) not null,  
    dept_name varchar(20),  
    salary     numeric(8,2))
```

- n **insert into** *instructor* **values** ('10211' , ' Smith' , ' Biology' , 66000);
- n **insert into** *instructor* **values** ('10211' , null, ' Biology' , 66000);



Integrity Constraints in Create Table

- n **not null**
- n **primary key** (A_1, \dots, A_n)
- n **foreign key** (A_m, \dots, A_n) **references** r

Example: Declare *ID* as the primary key for *instructor*

```
create table instructor (  
    ID          char(5),  
    name        varchar(20) not null,  
    dept_name   varchar(20),  
    salary      numeric(8,2),  
    primary key (ID),  
    foreign key (dept_name) references department)
```

primary key declaration on an attribute automatically ensures **not null**



And a Few More Relation Definitions

- n **create table** *student* (
 - ID* **varchar**(5),
 - name* **varchar**(20) not null,
 - dept_name* **varchar**(20),
 - tot_cred* **numeric**(3,0),
 - primary key** (*ID*),
 - foreign key** (*dept_name*) **references** *department*);

- n **create table** *takes* (
 - ID* **varchar**(5),
 - course_id* **varchar**(8),
 - sec_id* **varchar**(8),
 - semester* **varchar**(6),
 - year* **numeric**(4,0),
 - grade* **varchar**(2),
 - primary key** (*ID*, *course_id*, *sec_id*, *semester*, *year*),
 - foreign key** (*ID*) **references** *student*,
 - foreign key** (*course_id*, *sec_id*, *semester*, *year*) **references** *section*);

- | Note: *sec_id* can be dropped from primary key above, to ensure a student cannot be registered for two sections of the same course in the same semester



Even more

- n **create table** *course* (
 - course_id* **varchar(8) primary key,**
 - title* **varchar(50),**
 - dept_name* **varchar(20),**
 - credits* **numeric(2,0),**
 - foreign key** (*dept_name*) **references** *department*);
- | Primary key declaration can be combined with attribute declaration as shown above



Drop and Alter Table Constructs

n **drop table** *student*

- | Deletes the table and its contents

n **alter table**

| **alter table** *r* **add** *A* *D*

- ▶ where *A* is the name of the attribute to be added to relation *r* and *D* is the domain of *A*.
- ▶ All tuples in the relation are assigned *null* as the value for the new attribute.

| **alter table** *r* **drop** *A*

- ▶ where *A* is the name of an attribute of relation *r*
- ▶ Dropping of attributes not supported by many databases

| And more ...



Basic Query Structure

- n The SQL **data-manipulation language (DML)** provides the ability to query information, and insert, delete and update tuples
- n A typical SQL **query** has the form:

select A_1, A_2, \dots, A_n
from r_1, r_2, \dots, r_m
where P

- | A_i represents an attribute
 - | R_i represents a relation
 - | P is a predicate.
- n The result of an SQL query is a **relation**.



The select Clause

- n The **select** clause list the attributes desired in the result of a query
 - | corresponds to the projection operation of the relational algebra
- n Example: find the names of all instructors:

```
select name  
from instructor
```
- n NOTE: SQL keywords are case insensitive (i.e., you may use upper- or lower-case letters.)
 - | E.g. *Name* \equiv *NAME* \equiv *name*
 - | Some people use upper case wherever we use bold font.



The select Clause (Cont.)

- n SQL allows duplicates in relations as well as in query results.
- n To force the elimination of duplicates, insert the keyword **distinct** after select.
- n Find the names of all departments with instructor, and remove duplicates

```
select distinct dept_name  
from instructor
```

- n The (redundant) keyword **all** specifies that duplicates not be removed.

```
select all dept_name  
from instructor
```



The select Clause (Cont.)

- n An asterisk in the select clause denotes “all attributes”

```
select *  
from instructor
```

- n The **select** clause can contain arithmetic expressions involving the operation, +, −, *, and /, and operating on constants or attributes of tuples.
 - | Most systems also support additional functions
 - ▶ E.g., substring
 - | Most systems allow user defined functions (UDFs)
- n The query:

```
select ID, name, salary/12  
from instructor
```

would return a relation that is the same as the *instructor* relation, except that the value of the attribute *salary* is divided by 12.



The from Clause

- n The **from** clause lists the relations involved in the query
 - | Corresponds to the Cartesian product operation of the relational algebra.

- n Find the Cartesian product *instructor X teaches*

```
select *  
from instructor, teaches
```

- | generates every possible instructor – teaches pair, with all attributes from both relations
- n Cartesian product not very useful directly, but useful combined with where-clause condition (selection operation in relational algebra)



The where Clause

- n The **where** clause specifies conditions that the result must satisfy
 - | Corresponds to the selection predicate of the relational algebra.
- n To find all instructors in Comp. Sci. dept with salary > 80000

```
select name
from instructor
where dept_name = 'Comp. Sci.' and salary > 80000
```
- n Comparison results can be combined using the logical connectives **and**, **or**, and **not**.
- n Comparisons can be applied to results of arithmetic expressions.
- n SQL standard: any valid expression that returns a boolean result
 - | Vendor specific restrictions may apply!



Cartesian Product: *instructor X teaches*

instructor

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000

teaches

<i>ID</i>	<i>course_id</i>	<i>sec_id</i>	<i>semester</i>	<i>year</i>
10101	CS-101	1	Fall	2009
10101	CS-315	1	Spring	2010
10101	CS-347	1	Fall	2009
12121	FIN-201	1	Spring	2010
15151	MU-199	1	Spring	2010
22222	PHY-101	1	Fall	2009

<i>inst.ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>	<i>teaches.ID</i>	<i>course_id</i>	<i>sec_id</i>	<i>semester</i>	<i>year</i>
10101	Srinivasan	Comp. Sci.	65000	10101	CS-101	1	Fall	2009
10101	Srinivasan	Comp. Sci.	65000	10101	CS-315	1	Spring	2010
10101	Srinivasan	Comp. Sci.	65000	10101	CS-347	1	Fall	2009
10101	Srinivasan	Comp. Sci.	65000	12121	FIN-201	1	Spring	2010
10101	Srinivasan	Comp. Sci.	65000	15151	MU-199	1	Spring	2010
10101	Srinivasan	Comp. Sci.	65000	22222	PHY-101	1	Fall	2009
...
...
12121	Wu	Finance	90000	10101	CS-101	1	Fall	2009
12121	Wu	Finance	90000	10101	CS-315	1	Spring	2010
12121	Wu	Finance	90000	10101	CS-347	1	Fall	2009
12121	Wu	Finance	90000	12121	FIN-201	1	Spring	2010
12121	Wu	Finance	90000	15151	MU-199	1	Spring	2010
12121	Wu	Finance	90000	22222	PHY-101	1	Fall	2009
...
...



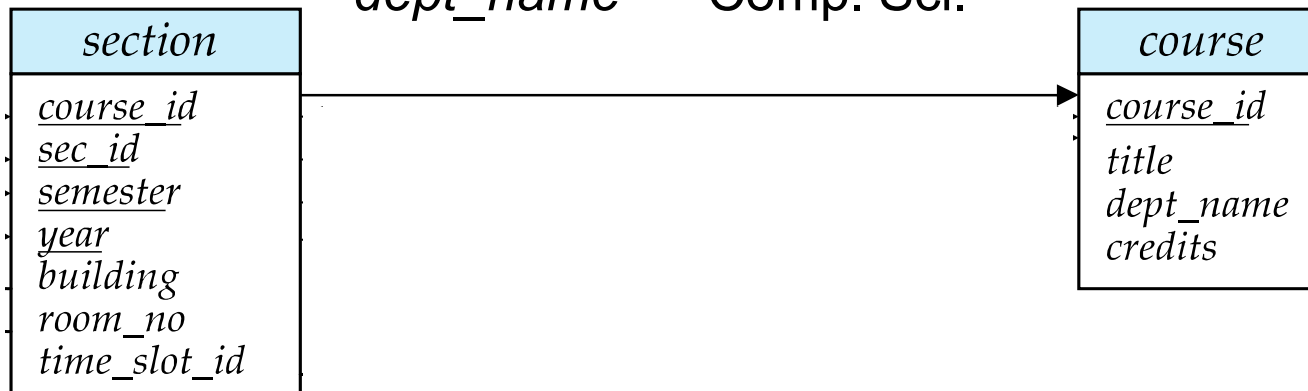
Joins

- n For all instructors who have taught some course, find their names and the course ID of the courses they taught.

```
select name, course_id  
from instructor, teaches  
where instructor.ID = teaches.ID
```

- n Find the course ID, semester, year and title of each course offered by the Comp. Sci. department

```
select section.course_id, semester, year, title  
from section, course  
where section.course_id = course.course_id and  
dept_name = 'Comp. Sci.'
```





Joined Relations

- n **Join operations** take two relations and return as a result another relation.
- n A join operation is a Cartesian product which requires that tuples in the two relations match (under some condition). It also specifies the attributes that are present in the result of the join
- n The join operations are typically used as subquery expressions in the **from** clause



Join operations – Example

n Relation *course*

<i>course_id</i>	<i>title</i>	<i>dept_name</i>	<i>credits</i>
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

n Relation *prereq*

<i>course_id</i>	<i>prereq_id</i>
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

n Observe that

prereq information is missing for CS-315 and
course information is missing for CS-437



Outer Join

- n An extension of the join operation that avoids loss of information.
- n Computes the join and then adds tuples from one relation that does not match tuples in the other relation to the result of the join.
- n Uses *null* values.



Joined Relations

- n **Join operations** take two relations and return as a result another relation.
- n These additional operations are typically used as subquery expressions in the **from** clause
- n **Join condition** – defines which tuples in the two relations match, and what attributes are present in the result of the join.
- n **Join type** – defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

<i>Join types</i>
inner join
left outer join
right outer join
full outer join

<i>Join Conditions</i>
natural
on <predicate>
using (A_1, A_1, \dots, A_n)



Joined Relations – Examples

- n *course* **inner join** *prereq* **on**
course.course_id = prereq.course_id

<i>course_id</i>	<i>title</i>	<i>dept_name</i>	<i>credits</i>	<i>prereq_id</i>	<i>course_id</i>
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190

- n What is the difference between the above, and a natural join?
- n *course* **left outer join** *prereq* **on**
course.course_id = prereq.course_id

<i>course_id</i>	<i>title</i>	<i>dept_name</i>	<i>credits</i>	<i>prereq_id</i>	<i>course_id</i>
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190
CS-315	Robotics	Comp. Sci.	3	<i>null</i>	<i>null</i>



Joined Relations – Examples

n **course natural right outer join prereq**

<i>course_id</i>	<i>title</i>	<i>dept_name</i>	<i>credits</i>	<i>prereq_id</i>
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	<i>null</i>	<i>null</i>	<i>null</i>	CS-101

n **course full outer join prereq using (course_id)**

<i>course_id</i>	<i>title</i>	<i>dept_name</i>	<i>credits</i>	<i>prereq_id</i>
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	<i>null</i>
CS-347	<i>null</i>	<i>null</i>	<i>null</i>	CS-101



The Rename Operation

- n The SQL allows renaming relations and attributes using the **as** clause:
old-name as new-name
- n E.g.
 - | **select** *ID, name, salary/12 as monthly_salary*
from *instructor*
- n Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci'.
 - | **select distinct** *T. name*
from *instructor as T, instructor as S*
where *T.salary > S.salary and S.dept_name = 'Comp. Sci.'*
- n Keyword **as** is optional and may be omitted
instructor as T \equiv *instructor T*
 - | Keyword **as** must be omitted in Oracle



String Operations

- n SQL includes a string-matching operator for comparisons on character strings. The operator “like” uses patterns that are described using two special characters:
 - | percent (%). The % character matches any substring.
 - | underscore (_). The _ character matches any character.
- n Find the names of all instructors whose name includes the substring “dar”.

```
select name  
from instructor  
where name like '%dar%'
```

- n Match the string “100 %”

```
like '100 \%' escape '\'
```



String Operations (Cont.)

- n Patterns are case sensitive.
- n Pattern matching examples:
 - | 'Intro%' matches any string beginning with "Intro".
 - | '%Comp%' matches any string containing "Comp" as a substring.
 - | '___' matches any string of exactly three characters.
 - | '___ %' matches any string of at least three characters.
- n SQL supports a variety of string operations such as
 - | concatenation (using "||")
 - | converting from upper to lower case (and vice versa)
 - | finding string length, extracting substrings, etc.



Case Construct

- n Like case, if, and ? Operators in programming languages

case

when c_1 **then** e_1

when c_2 **then** e_2

...

[else e_n **]**

end

- n Each c_i is a condition
- n Each e_i is an expression
- n Returns the first e_i for which c_i evaluates to *true*
 - | If none of the c_i is true, then return e_n (**else**)
 - ▶ If there is no else return *null*



Case Construct Example

- n Like case, if, and ? Operators in programming languages

select

name,

case

when salary > 1000000 **then** 'premium'

else 'standard'

end as customer_group

from customer



Ordering the Display of Tuples

- n List in alphabetic order the names of all instructors
select distinct *name*
from *instructor*
order by *name*
- n We may specify **desc** for descending order or **asc** for ascending order, for each attribute; ascending order is the default.
 - | Example: **order by** *name desc*
- n Can sort on multiple attributes
 - | Example: **order by** *dept_name, name*
- n **Order is not expressible in the relational model!**



Where Clause Predicates

- n SQL includes a **between** comparison operator
- n Example: Find the names of all instructors with salary between \$90,000 and \$100,000 (that is, \geq \$90,000 and \leq \$100,000)
 - | **select** *name*
from *instructor*
where *salary* **between** 90000 **and** 100000
- n Tuple comparison
 - | **select** *name, course_id*
from *instructor, teaches*
where (*instructor.ID, dept_name*) = (*teaches.ID, 'Biology'*);



Set Operations

- n Find courses that ran in Fall 2009 or in Spring 2010

(select *course_id* from *section* where *sem* = 'Fall' and *year* = 2009)

union

(select *course_id* from *section* where *sem* = 'Spring' and *year* = 2010)

- n Find courses that ran in Fall 2009 and in Spring 2010

(select *course_id* from *section* where *sem* = 'Fall' and *year* = 2009)

intersect

(select *course_id* from *section* where *sem* = 'Spring' and *year* = 2010)

- n Find courses that ran in Fall 2009 but not in Spring 2010

(select *course_id* from *section* where *sem* = 'Fall' and *year* = 2009)

except

(select *course_id* from *section* where *sem* = 'Spring' and *year* = 2010)



Set Operations

- n Set operations **union**, **intersect**, and **except**
 - | Each of the above operations automatically eliminates duplicates
- n To retain all duplicates use the corresponding multiset versions **union all**, **intersect all** and **except all**.

Suppose a tuple occurs m times in r and n times in s , then, it occurs:

- | $m + n$ times in r **union all** s
- | $\min(m, n)$ times in r **intersect all** s
- | $\max(0, m - n)$ times in r **except all** s



Null Values

- n It is possible for tuples to have a null value, denoted by *null*, for some of their attributes
- n *null* signifies an unknown value or that a value does not exist.
- n The result of any arithmetic expression and comparisons involving *null* evaluate to *null*
 - | Example: $5 + \textit{null}$ returns *null*
 $\textit{null} > 5$ returns *null*
 $\textit{null} = \textit{null}$ returns *null*
- n The predicate **is null** can be used to check for null values.
 - | Example: Find all instructors whose salary is null.
select *name*
from *instructor*
where *salary is null*



Null Values and Three Valued Logic

- n Any comparison with *null* returns *null*
 - | Example: $5 < null$ or $null \neq null$ or $null = null$
- n Three-valued logic using the truth value *null*:
 - | OR: $(null \text{ or } true) = true$,
 $(null \text{ or } false) = null$
 $(null \text{ or } null) = null$
 - | AND: $(true \text{ and } null) = null$,
 $(false \text{ and } null) = false$,
 $(null \text{ and } null) = null$
 - | NOT: $(\text{not } null) = null$
 - | “*P* is null” evaluates to true if predicate *P* evaluates to *null*
- n Result of **where** clause predicate is treated as *false* if it evaluates to *null*



Aggregate Functions

- n These functions operate on the multiset of values of a column of a relation, and return a value
 - avg:** average value
 - min:** minimum value
 - max:** maximum value
 - sum:** sum of values
 - count:** number of values

- n Most DBMS support user defined aggregation functions



Aggregate Functions (Cont.)

- n Find the average salary of instructors in the Computer Science department
 - | **select avg** (*salary*)
from *instructor*
where *dept_name*= ' Comp. Sci.' ;
- n Find the total number of instructors who teach a course in the Spring 2010 semester
 - | **select count** (**distinct** *ID*)
from *teaches*
where *semester* = ' Spring' **and** *year* = 2010
- n Find the number of tuples in the *course* relation
 - | **select count** (*)
from *course*;



Aggregate Functions – Group By

- n Find the average salary of instructors in each department
 - | **select** *dept_name*, **avg** (*salary*)
from *instructor*
group by *dept_name*;
 - | Note: departments with no instructor will not appear in result

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000

<i>dept_name</i>	<i>avg_salary</i>
Biology	72000
Comp. Sci.	77333
Elec. Eng.	80000
Finance	85000
History	61000
Music	40000
Physics	91000



Aggregation (Cont.)

- n Attributes in **select** clause outside of aggregate functions must appear in **group by** list
 - | */* erroneous query */*
select *dept_name*, **ID**, **avg** (*salary*)
from *instructor*
group by *dept_name*;



Aggregate Functions – Having Clause

- n Find the names and average salaries of all departments whose average salary is greater than 42000

```
select dept_name, avg (salary)
from instructor
group by dept_name
having avg (salary) > 42000;
```

Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups



Null Values and Aggregates

- n Total all salaries

```
select sum (salary )  
from instructor
```

- | Above statement ignores null amounts
- | Result is *null* if there is no non-null amount
- n All aggregate operations except **count(*)** ignore tuples with null values on the aggregated attributes
- n What if collection has only null values?
 - | count returns 0
 - | all other aggregates return null



Empty Relations and Aggregates

- n What if the input relation is empty
- n Conventions:
 - | **sum**: returns *null*
 - | **avg**: returns *null*
 - | **min**: returns *null*
 - | **max**: returns *null*
 - | **count**: returns 0



Duplicates

- n In relations with duplicates, SQL can define how many copies of tuples appear in the result.
- n **Multiset (bag semantics)** versions of some of the relational algebra operators – given multiset relations r_1 and r_2 :
 1. $\sigma_{\theta}(r_1)$: If there are c_1 copies of tuple t_1 in r_1 , and t_1 satisfies selections σ_{θ} , then there are c_1 copies of t_1 in $\sigma_{\theta}(r_1)$.
 2. $\Pi_A(r)$: For each copy of tuple t_1 in r_1 , there is a copy of tuple $\Pi_A(t_1)$ in $\Pi_A(r_1)$ where $\Pi_A(t_1)$ denotes the projection of the single tuple t_1 .
 3. $r_1 \times r_2$: If there are c_1 copies of tuple t_1 in r_1 and c_2 copies of tuple t_2 in r_2 , there are $c_1 \times c_2$ copies of the tuple $t_1 \cdot t_2$ in $r_1 \times r_2$



Multiset Relational Algebra

- n Pure relational algebra operates on **set-semantics** (no duplicates allowed)
 - | e.g. after projection
- n Multiset (**bag-semantics**) relational algebra retains duplicates, to match SQL semantics
 - | SQL duplicate retention was initially for efficiency, but is now a feature
- n Multiset relational algebra defined as follows
 - | **selection**: has as many duplicates of a tuple as in the input, if the tuple satisfies the selection
 - | **projection**: one tuple per input tuple, even if it is a duplicate
 - | **cross product**: If there are m copies of $t1$ in r , and n copies of $t2$ in s , there are $m \times n$ copies of $t1.t2$ in $r \times s$
 - | Other operators similarly defined
 - ▶ E.g. **union**: $m + n$ copies, **intersection**: $\min(m, n)$ copies
difference: $\max(0, m - n)$ copies



Duplicates (Cont.)

- n Example: Suppose multiset relations $r_1 (A, B)$ and $r_2 (C)$ are as follows:

$$r_1 = \{(1, a) (2, a)\} \quad r_2 = \{(2), (3), (3)\}$$

- n Then $\Pi_B(r_1)$ would be $\{(a), (a)\}$, while $\Pi_B(r_1) \times r_2$ would be $\{(a,2), (a,2), (a,3), (a,3), (a,3), (a,3)\}$

- n SQL duplicate semantics:

```
select  $A_1, A_2, \dots, A_n$   
from  $r_1, r_2, \dots, r_m$   
where  $P$ 
```

is equivalent to the *multiset* version of the expression:

$$\square_{A_1, A_2, \dots, A_n} (\sigma_P (r_1 \square r_2 \square \dots \square r_m))$$



SQL and Relational Algebra

n **select** A_1, A_2, \dots, A_n
from r_1, r_2, \dots, r_m
where P

is equivalent to the following expression in multiset relational algebra

$$\Pi_{A_1, \dots, A_n} (\sigma_P (r_1 \times r_2 \times \dots \times r_m))$$

n **select** $A_1, A_2, \text{sum}(A_3)$
from r_1, r_2, \dots, r_m
where P
group by A_1, A_2

is equivalent to the following expression in multiset relational algebra

$$A_1, A_2 \mathcal{G} \text{sum}(A_3) (\sigma_P (r_1 \times r_2 \times \dots \times r_m))$$



SQL and Relational Algebra

- n More generally, the non-aggregated attributes in the **select** clause may be a subset of the **group by** attributes, in which case the equivalence is as follows:

```
select  $A_1, \text{sum}(A_3)$  AS  $\text{sumA3}$   
from  $r_1, r_2, \dots, r_m$   
where P  
group by  $A_1, A_2$ 
```

is equivalent to the following expression in multiset relational algebra

$$\Pi_{A_1, \text{sumA3}}(\rho_{A_1, A_2} \mathcal{G} \text{sum}(A_3) \text{ as } \text{sumA3}(\sigma_P(r_1 \times r_2 \times \dots \times r_m)))$$



Subqueries in the From Clause

- n SQL allows a subquery expression to be used in the **from** clause
- n Find the average instructors' salaries of those departments where the average salary is greater than \$42,000.

```
select dept_name, avg_salary
from (select dept_name, avg (salary) as avg_salary
      from instructor
      group by dept_name)
where avg_salary > 42000;
```

- n Note that we do not need to use the **having** clause
- n Another way to write above query

```
select dept_name, avg_salary
from (select dept_name, avg (salary)
      from instructor
      group by dept_name)
      as dept_avg (dept_name, avg_salary)
where avg_salary > 42000;
```




Nested Subqueries

- n SQL provides a mechanism for the nesting of subqueries.
- n A **subquery** is a **select-from-where** expression that is nested within another query.
- n A common use of subqueries is to perform tests for set membership, set comparisons, and set cardinality.



Example Query

- n Find courses offered in Fall 2009 and in Spring 2010

```
select distinct course_id
from section
where semester = ' Fall' and year= 2009 and
       course_id in (select course_id
                       from section
                       where semester = ' Spring' and year= 2010);
```

- n Find courses offered in Fall 2009 but not in Spring 2010

```
select distinct course_id
from section
where semester = ' Fall' and year= 2009 and
       course_id not in (select course_id
                             from section
                             where semester = ' Spring' and year=
2010);
```



Example Query

- n Find the total number of (distinct) students who have taken course sections taught by the instructor with *ID* 10101

```
select count (distinct ID)  
from takes  
where (course_id, sec_id, semester, year) in  
        (select course_id, sec_id, semester, year  
         from teaches  
         where teaches.ID= 10101);
```

- n **Note:** Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features.



Quantification

- n Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

```
select distinct T.name  
from instructor as T, instructor as S  
where T.salary > S.salary and S.dept_name = ' Biology' ;
```

- n Same query using > **some** clause

```
select name  
from instructor  
where salary > some (select salary  
                        from instructor  
                        where dept_name = ' Biology' );
```



Definition of Some Clause

n $F <comp> \mathbf{some} r \Leftrightarrow \exists t \in r \text{ such that } (F <comp> t)$

Where $<comp>$ can be: $<, \leq, >, =, \neq$

$(5 < \mathbf{some} \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{true}$ (read: 5 < some tuple in the relation)

$(5 < \mathbf{some} \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{false}$

$(5 = \mathbf{some} \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$

$(5 \neq \mathbf{some} \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true (since } 0 \neq 5)$

$(= \mathbf{some}) \equiv \mathbf{in}$

However, $(\neq \mathbf{some}) \not\equiv \mathbf{not in}$



Example Query

- n Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

```
select name  
from instructor  
where salary > all (select salary  
                        from instructor  
                        where dept_name = ' Biology' );
```



Definition of all Clause

n $F \text{ <comp> all } r \Leftrightarrow \forall t \in r (F \text{ <comp> } t)$

$(5 < \text{all } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$

$(5 < \text{all } \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$

$(5 = \text{all } \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$

$(5 \neq \text{all } \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$

$(\neq \text{all}) \equiv \text{not in}$

However, $(= \text{all}) \not\equiv \text{in}$



Test for Empty Relations

- n The **exists** construct returns the value **true** if the argument subquery returns a nonempty result.
- n **exists** $r \Leftrightarrow r \neq \emptyset$
- n **not exists** $r \Leftrightarrow r = \emptyset$



Correlation Variables

- n Yet another way of specifying the query “Find all courses taught in both the Fall 2009 semester and in the Spring 2010 semester”

```
select course_id
from section as S
where semester = ' Fall' and year= 2009 and
exists (select *
         from section as T
         where semester = ' Spring' and year= 2010
         and S.course_id= T.course_id);
```

- n **Correlated subquery**
- n **Correlation name** or **correlation variable**



Not Exists

- n Find all students who have taken all courses offered in the Biology department.

```
select distinct S.ID, S.name
from student as S
where not exists ( (select course_id
                    from course
                    where dept_name = ' Biology' )
except
                    (select T.course_id
                    from takes as T
                    where S.ID = T.ID));
```

- n Note that $X - Y = \emptyset \Leftrightarrow X \subseteq Y$
- n *Note:* Cannot write this query using = **all** and its variants



Test for Absence of Duplicate Tuples

- n The **unique** construct tests whether a subquery has any duplicate tuples in its result.
 - | (Evaluates to “true” on an empty set)
- n Find all courses that were offered at most once in 2009

```
select T.course_id
from course as T
where unique (select R.course_id
                 from section as R
                 where T.course_id= R.course_id
                 and R.year = 2009);
```



With Clause

- n The **with** clause provides a way of defining a temporary view whose definition is available only to the query in which the **with** clause occurs.
- n Find all departments with the maximum budget

```
with max_budget (value) as  
    (select max(budget)  
     from department)  
select budget  
from department, max_budget  
where department.budget = max_budget.value;
```



Scalar Subquery

- n Scalar subquery is one which is used where a single value is expected
- n E.g.

```
select dept_name,  
      (select count(*)  
       from instructor  
       where department.dept_name = instructor.dept_name)  
      as num_instructors  
from department;
```
- n E.g.

```
select name  
from instructor  
where salary * 10 >  
      (select budget from department  
       where department.dept_name = instructor.dept_name)
```
- n Runtime error if subquery returns more than one result tuple



Query Features Recap - Syntax

n An SQL query is either a Select-from-where block or a set operation

n An SQL query block is structured like this:

SELECT [**DISTINCT**] select_list

[**FROM** from_list]

[**WHERE** where_condition]

[**GROUP BY** group_by_list]

[**HAVING** having_condition]

[**ORDER BY** order_by_list]

n Set operations

[Query Block] set_op [Query Block]

set_op: [**ALL**] **UNION** | **INTERSECT** | **EXCEPT**



Query Features Recap - Syntax

n Almost all clauses are optional

n Examples:

- | **SELECT * FROM r;**

- | **SELECT 1;**

 - ▶ Convention: returns single tuple

- | **SELECT 'ok' FROM accounts HAVING sum(balance) = 0;**

- | **SELECT 1 GROUP BY 1;**

- | **SELECT 1 HAVING true;**

- | Let r be a relation with two attributes a and b

 - ▶ **SELECT a,b FROM r**

 - WHERE a IN (SELECT a FROM r) AND b IN (SELECT b FROM r)**

 - GROUP BY a,b HAVING count(*) > 0;**

n Note:

- | Not all systems support all of this “non-sense”



Syntax - SELECT

- n **SELECT** [**DISTINCT** [**ON** (*distinct_list*)]] *select_list*
- n *select_list*
 - | List of projection expressions
 - ▶ [*expr*] [**AS** *name*]
 - | *expr*
 - ▶ Expression over attributes, constants, arithmetic operators, functions, **CASE**-construct, aggregation functions
- n *distinct_list*
 - | List of expressions
- n Examples:
 - | **SELECT DISTINCT ON** (*a % 2*) *a* **FROM** *r*;
 - | **SELECT** *substring(a, 1,2)* **AS** *x* **FROM** *r*;
 - | **SELECT CASE WHEN** *a = 2* **THEN** *a* **ELSE** *null* **END AS** *b* **FROM** *r*;
 - | **SELECT** *a = b* **AS** *is_a_equal_to_b* **FROM** *r*;



Syntax - FROM

- n **FROM** from_list
- n from_list
 - | List of from clause expressions
 - ▶ subquery | relation | constant_relation | join_expr [alias]
 - | subquery
 - ▶ Any valid SQL query – alias is not optional
 - | relation
 - ▶ A relation in the database
 - | constant_relation
 - ▶ (**VALUES** tuples) – alias is not optional
 - | join_expr
 - ▶ joins between from_clause entries
 - | alias
 - ▶ [AS] b [(attribute_name_list)]



Syntax – FROM (cont.)

- n Examples (relation r with attributes a and b):
 - | **SELECT * FROM r;**
 - | **SELECT * FROM r AS g(v,w);**
 - | **SELECT * FROM r x;**
 - | **SELECT * FROM (VALUES (1,2), (3,1)) AS s(u,v);**
 - | **SELECT * FROM r NATURAL JOIN s, t;**
 - | **SELECT * FROM ((r JOIN s ON (r.a = s.c)) NATURAL JOIN (SELECT * FROM t) AS new);**
 - | **SELECT * FROM (SELECT * FROM r) AS r;**
 - | **SELECT * FROM (SELECT * FROM (SELECT * FROM r) AS r) AS r;**



Syntax - WHERE

- n **WHERE** where_condition
- n where_condition: A boolean expression over
 - | Attributes
 - | Constants: e.g., true, 1, 0.5, 'hello'
 - | Comparison operators: =, <, >, **IS DISTINCT FROM**, **IS NULL**, ...
 - | Arithmetic operators: +, -, /, %
 - | Function calls
 - | Nested subquery expressions
- n Examples
 - | **SELECT * FROM r WHERE a = 2;**
 - | **SELECT * FROM r WHERE true OR false;**
 - | **SELECT * FROM r WHERE NOT(a = 2 OR a = 3);**
 - | **SELECT * FROM r WHERE a IS DISTINCT FROM b;**
 - | **SELECT * FROM r WHERE a < ANY (SELECT c FROM s);**
 - | **SELECT * FROM r WHERE a = (SELECT count(*) FROM s);**



Syntax – GROUP BY

- n **GROUP BY** group_by_list
- n group_by_list
 - | List of expressions
 - ▶ Expression over attributes, constants, arithmetic operators, functions, **CASE**-construct, aggregation functions
- n Examples:
 - | **SELECT** sum(a), b **FROM** r **GROUP BY** b;
 - | **SELECT** sum(a), b, c **FROM** r **GROUP BY** b, c;
 - | **SELECT** sum(a), b/2 **FROM** r **GROUP BY** b/2;
 - | **SELECT** sum(a), b **FROM** r **GROUP BY** b > 5;
 - ▶ Incorrect, cannot select b, because it is not an expression in the group by clause
 - | **SELECT** sum(a), b **FROM** r **GROUP BY** b **IN** (**SELECT** c **FROM** s);



Syntax – HAVING

- n **HAVING** having_condition
- n having_condition
 - | Like where_condition except that expressions over attributes have either to be in the **GROUP BY** clause or are aggregated
- n Examples:
 - | **SELECT** sum(a), b **FROM** r **GROUP BY** b **HAVING** sum(a) > 10;
 - | **SELECT** sum(a), b **FROM** r **GROUP BY** b **HAVING** sum(a) + 5 > 10;
 - | **SELECT** sum(a), b **FROM** r **GROUP BY** b **HAVING** true;
 - | **SELECT** sum(a), b **FROM** r **GROUP BY** b **HAVING** count(*) = 50;
 - | **SELECT** b **FROM** r **GROUP BY** b **HAVING** sum(a) > 10;



Syntax – ORDER BY

n **ORDER BY** order_by_list

n order_by_list

| Like select_list minus renaming

| Optional [**ASC** | **DESC**] for each item

n Examples:

| **SELECT * FROM r ORDER BY a;**

| **SELECT * FROM r ORDER BY b, a;**

| **SELECT * FROM r ORDER BY a * 2;**

| **SELECT * FROM r ORDER BY a * 2, a;**

| **SELECT * FROM r ORDER BY a + (SELECT count(*) FROM s);**



Query Semantics

- n Evaluation Algorithm (you can do it manually – sort of)
 - 1. Compute **FROM** clause
 - 1. Compute cross product of all items in the **FROM** clause
 - ▶ Relations: nothing to do
 - ▶ Subqueries: use this algorithm to recursively compute the result of subqueries first
 - ▶ Join expressions: compute the join
 - 2. Compute **WHERE** clause
 - 1. For each tuple in the result of 1. evaluate the **WHERE** clause condition
 - 3. Compute **GROUP BY** clause
 - 1. Group the results of step 2. on the **GROUP BY** expressions
 - 4. Compute **HAVING** clause
 - 1. For each group (if any) evaluate the **HAVING** condition



Query Semantics (Cont.)

5. Compute **ORDER BY** clause
 5. Order the result of step 4 on the **ORDER BY** expressions
6. Compute **SELECT** clause
 5. Project each result tuple from step 5 on the **SELECT** expressions
- n If the **WHERE**, **SELECT**, **GROUP BY**, **HAVING**, **ORDER BY** clauses have any nested subqueries
 - I For each tuple t in the result of the **FROM** clause
 - ▶ Substitute the correlated attributes with values from t
 - ▶ Evaluate the resulting query
 - ▶ Use the result to evaluate the expression in the clause the subquery occurs in



Query Semantics (Cont.)

- n Equivalent relational algebra expression
 - | **ORDER BY** has no equivalent, because relations are unordered
 - | Nested subqueries: need to extend algebra (not covered here)
- n Each query block is equivalent to

$$\pi(\sigma(\mathcal{G}(\pi(\sigma(F_1 \times \dots \times F_n))))))$$

- n Where F_i is the translation of the i^{th} **FROM** clause item
- n Note: we leave out the arguments



Modification of the Database

- n Deletion of tuples from a given relation
- n Insertion of new tuples into a given relation
- n Updating values in some tuples in a given relation



Modification of the Database – Deletion

- n Delete all instructors

```
delete from instructor
```

- n Delete all instructors from the Finance department

```
delete from instructor  
where dept_name = ' Finance' ;
```

- n Delete all tuples in the *instructor* relation for those instructors associated with a department located in the Watson building.

```
delete from instructor  
where dept_name in (select dept_name  
                        from department  
                        where building = ' Watson' );
```



Deletion (Cont.)

- n Delete all instructors whose salary is less than the average salary of instructors

```
delete from instructor  
where salary < (select avg (salary) from instructor);
```

- | Problem: as we delete tuples from instructor, the average salary changes
- | Solution used in SQL:
 1. First, compute **avg** salary and find all tuples to delete
 2. Next, delete all tuples found above (without recomputing **avg** or retesting the tuples)



Modification of the Database – Insertion

- n Add a new tuple to *course*

```
insert into course  
values ( ' CS-437' , ' Database Systems' , ' Comp. Sci.' , 4);
```

- n or equivalently

```
insert into course (course_id, title, dept_name, credits)  
values ( ' CS-437' , ' Database Systems' , ' Comp. Sci.' , 4);
```

- n Add a new tuple to *student* with *tot_creds* set to null

```
insert into student  
values ( ' 3003' , ' Green' , ' Finance' , null);
```



Insertion (Cont.)

- n Add all instructors to the *student* relation with *tot_creds* set to 0

```
insert into student
```

```
  select ID, name, dept_name, 0
```

```
  from instructor
```

- n The **select from where** statement is evaluated fully before any of its results are inserted into the relation (otherwise queries like

```
  insert into table1 select * from table1
```

would cause problems, if *table1* did not have any primary key defined.



Modification of the Database – Updates

n Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others receive a 5% raise

| Write two **update** statements:

```
update instructor
```

```
  set salary = salary * 1.03
```

```
  where salary > 100000;
```

```
update instructor
```

```
  set salary = salary * 1.05
```

```
  where salary <= 100000;
```

| The order is important

| Can be done better using the **case** statement (next slide)



Case Statement for Conditional Updates

- n Same query as before but with case statement

```
update instructor  
  set salary = case  
    when salary <= 100000 then salary * 1.05  
    else salary * 1.03  
  end
```




Recap

- n SQL queries
 - | Clauses: **SELECT**, **FROM** , **WHERE**, **GROUP BY**, **HAVING**, **ORDER BY**
 - | Nested subqueries
 - | Equivalence with relational algebra
- n SQL update, inserts, deletes
 - | Semantics of referencing updated relation in **WHERE**
- n SQL DDL
 - | Table definition: **CREATE TABLE**



End of Chapter 4

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