

CS440: Programming Languages and Translators

Lecture 11

Spring 2023

Environments with functions: first try

```
let add = fun (x, y) -> x + y  
let three = add (1, 2)
```

Var	Value

Environments with functions: first try

```
let add = fun (x, y) -> x + y  
let three = add (1, 2)
```

Var	Value
add	fun (x, y) -> x + y

Environments with functions: first try

```
let add = fun (x, y) -> x + y  
let three = add (1, 2)
```

Var	Value
add	fun (x, y) -> x + y
x	1
y	2

Environments with functions: first try

```
let add = fun (x, y) -> x + y  
let three = add (1, 2)
```

Var	Value
add	fun (x, y) -> x + y
three	3

Environments with functions: first try

```
let add = fun x -> fun y -> x + y
let add1 = add 1
let three = add1 2
```

Var	Value
add	fun x -> fun y -> x + y

Environments with functions: first try

```
let add = fun x -> fun y -> x + y
let add1 = add 1
let three = add1 2
```

Var	Value
add	fun x -> fun y -> x + y
x	1

Environments with functions: first try

```
let add = fun x -> fun y -> x + y
let add1 = add 1
let three = add1 2
```

Var	Value
add	fun x -> fun y -> x + y
add1	fun y -> x + y

Uh oh

Environments with functions: first try

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value

Environments with functions: first try

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value
x	1

Environments with functions: first try

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value
x	1
f	fun y -> x + y

Environments with functions: first try

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value
x	1
f	fun y -> x + y
x	2

Environments with functions: first try

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value
x	1
f	fun y -> x + y
x	2
y	2

x should still be 1 in f!

Second try: use *closures*

- Closure: function code + environment
- This will be the value of a function

- in ML:

`type value = ...`

- `| VClos of var * expr * env`
`and env = (var * value option ref) list`



Argument name

With closures

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value

With closures

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value
x	1

With closures

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Var	Value				
x	1				
f	(fun y -> x + y, <table border="1"><thead><tr><th>Var</th><th>Value</th></tr></thead><tbody><tr><td>x</td><td>1</td></tr></tbody></table>)	Var	Value	x	1
Var	Value				
x	1				

With closures

```
let x = 1 in
let f y = x + y in
let x = 2 in
f 2
```

Var	Value				
x	1				
f	(fun y -> x + y, <table border="1"><thead><tr><th>Var</th><th>Value</th></tr></thead><tbody><tr><td>x</td><td>1</td></tr></tbody></table>)	Var	Value	x	1
Var	Value				
x	1				
x	2				

With closures

```
let x = 1 in  
let f y = x + y in  
let x = 2 in  
f 2
```

Call the function with the
environment from the closure
(+ arguments)

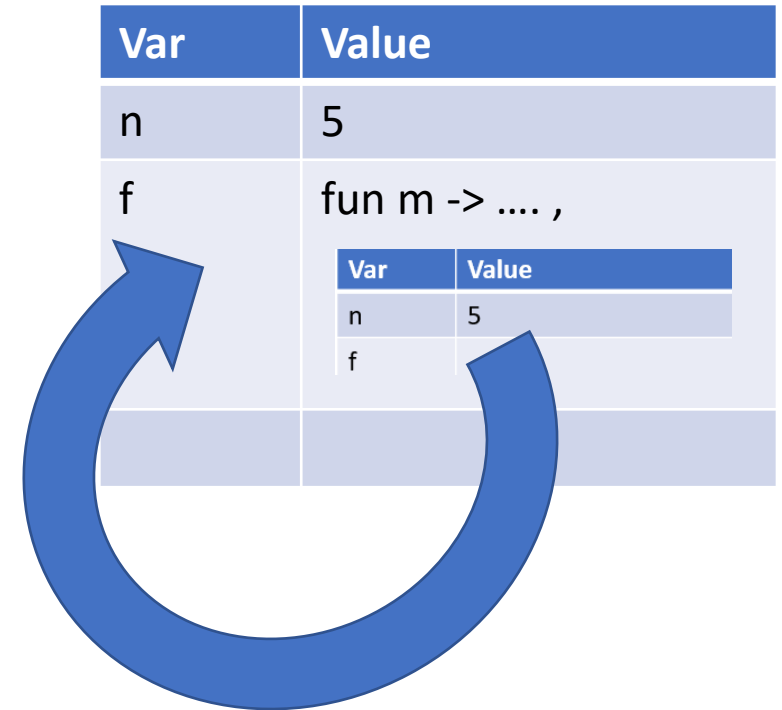
Var	Value
x	1
y	2

Interpreting with closures

- Interpreting a function: `fun x -> e`
 - Return a closure with variable `x`, expression `e`, current environment
- Interpreting an application `e1 e2`
 - Interpret `e1` to closure `(x, e, env)`
 - Interpret `e2` to arg value `v`
 - Add `x -> v` to `env`, interpret `e` with this `env`

Recursive closures

```
let n = 5 in
let rec f m =
  if m >= n then 1
  else m * f (m + 1)
in
f 0
```



Interpreting with recursive closures

- Interpreting a function def: `let rec f x = e1 in e2`
 - Let env' = current env extended with placeholder for f
 - Let $clos = \lambda clos(x, e1, env')$
 - Update env' with $f \rightarrow clos$
 - Evaluate $e2$ with env'
- Interpreting an application `e1 e2`
 - Same as before: when you evaluate $e1$ to a closure, the function is already in the environment