

Name

CWID

# Midterm Exam

March 12th, 2020

10:00-11:15

## CS520 - Data Integration, Warehousing, and Provenance

### Results

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*Please leave this empty!*

1.1

1.2

1.3

Sum

# Instructions

- Try to answer all the questions using what you have learned in class. Keep hard questions until the end.
- **When writing a query, write the query in a way that it would work over all possible database instances and not just for the given example instance!**
- The exam is closed book and closed notes! No calculator, smartphones, or similar allowed!

Consider the following database schema and example instance about music albums:

## user

nickname	name	post Visibility	country
BobAwesome	Bob	FOF	USA
Ali12	Alice	friends	France
Peter	Peter	friends	India
Pokegert	Gert	public	China

## friends

person	friend
BobAwesome	Ali12
Ali12	BobAwesome
BobAwesome	Peter
Peter	Pokegert

## posts

pid	user	text	time
1	BobAwesome	Hello just brought ...	2018-01-10
2	BobAwesome	meet @Ali12 at ...	2018-01-11
3	Peter	... is great, would recommend.	2018-01-15

### Hints:

- Attributes with black background form the primary key of a relation (e.g., *nickname* for relation *user*)
- The attributes *person* and *friend* of relation *friends* are a foreign key to relation *user*.
- The attribute *user* of relation *posts* is a foreign key to relation *user*.

## Part 1.1 Datalog (Total: 38 Points)

Recall that Datalog applies set semantics.

### Question 1.1.1 (5 Points)

Write a **Datalog program** that returns the *name* and *nickname* of users from USA.

### Solution

```
Q(NN,N) :- user(N,N,_,usa).
```

### Question 1.1.2 (7 Points)

Write a **Datalog program** that returns the *names* of users which are living in France and have posted before 2018-01-10 or live in USA and whose posts are visible to the public (post visibility = public).

### Solution

```
Q(X) :- user(N,X,_,france), posts(_,N,_,T), T < 2018-01-10.  
Q(X) :- user(_,X,public,usa).
```

### Question 1.1.3 (8 Points)

Write a **Datalog program** that returns the names of users that are not friends of user *BobAwesome* nor are they friends of a friend of *BobAwesome*. For example, in the example EDB instance there is no such person.

### Solution

```
friendsOfBob(X) :- friends(bobawesome,X).
friendsOfBob(X) :- friends(bobawesome,Y), friends(Y,X).
Q(X) :- users(X,_,_,_), not friendsOfBob(X).
```

### Question 1.1.4 (9 Points)

Write a **Datalog program** that returns users who have friends or friends of friends in every country.

### Solution

```
fofCntr(X,C) :- friends(X,Y), user(Y,_,_,C).
fofCntr(X,C) :- friends(X,Z), friends(Z,Y), user(Y,_,_,C).
countries(X) :- user(_,_,_,X).
U(X) :- user(X,_,_,_).
missingCountry(X) :- , country(C), not fofCntr(X,C).
Q(X) :- U(X), not missingCountry(X).
```

### Question 1.1.5 (9 Points)

Write a **Datalog program** that returns pairs of countries (C1,C2) such that there exists at least one path in the friendship graph that connects a user from country C1 with a user from country C2. **Here we do not care about the direction of edges, e.g., there is a path from Pokegert to Ali12.**

### Solution

```
TC(X,Y) :- friends(X,Y).
TC(X,Y) :- friends(Y,X).
TC(X,Y) :- TC(X,Z), TC(Z,Y).
Q(C1,C2) :- user(X,_,_,C1), user(Y,_,_,C2), TC(X,Y).
```

## Part 1.2 Constraints (Total: 26 Points)

### Question 1.2.1 Expressing Constraints in First-Order Logic (13 Points)

Recall the representation of constraints as universally quantified formulas in first-order logic introduced in class. Write down the logical encoding of the following constraints over the example schema:

- The foreign key from attribute `friend` of relation `friends` to relation `user`.
- Friendship has to be reciprocal, i.e., if `X` is a friend of `Y`, then also `Y` has to be a friend of `X`.
- The primary key of relation `posts`
- The following functional dependency for relation `users`:  $country \rightarrow postVisibility$

### Solution

$$FK_1 : \forall p, f : friends(p, f) \rightarrow \exists x_1, x_2, x_3 : user(f, x_1, x_2, x_3)$$

$$REPROC : \forall x, y : friends(x, y) \rightarrow friends(y, x)$$

$$PK : \forall p, u_1, t_1, i_1, p, f_2, t_2, i_2 : posts(p, u_1, t_1, i_1) \wedge posts(p, u_2, t_2, i_2) \rightarrow u_1 = u_2 \wedge t_1 = t_2 \wedge i_1 = i_2$$

$$FD : \forall n_1, a_1, p_1, c, n_2, a_2, p_2 : user(n_1, a_1, p_1, c), user(n_2, a_2, p_2, c) \rightarrow p_1 = p_2$$

## Question 1.2.2 Creating Denial Constraints (13 Points)

Create denial constraints over the example schema based on the following descriptions.

- The friendship graph is not allowed to contain any triangles, i.e., this constraint is violated if there exists users  $X$ ,  $Y$ , and  $Z$  such that  $X$  is a friend of  $Y$ ,  $Y$  is a friend of  $Z$ , and  $Z$  is a friend of  $X$  (Note that the direction of edges matters!).
- Users of country USA are not allowed to post after 2020-03-11.
- Implement the primary key of relation users.

## Solution

$$d_1 : \forall u_1, u_2, u_3 : \neg(\text{friends}(u_1, u_2) \wedge \text{friends}(u_2, u_3) \wedge \text{friends}(u_3, u_1))$$

$$d_2 : \forall n, a, p, e, t : \neg(\text{user}(n, a, p, \text{USA}) \wedge \text{posts}(p, n, e, t) \wedge t > 2020 - 03 - 11)$$

$$d_3 : \forall u, n_1, p_1, c_1, n_2, p_2, c_2 : \neg(\text{user}(u, n_1, p_1, c_1) \wedge \text{user}(u, n_2, p_2, c_2) \wedge n_1 \neq n_2)$$

$$\forall u, n_1, p_1, c_1, n_2, p_2, c_2 : \neg(\text{user}(u, n_1, p_1, c_1) \wedge \text{user}(u, n_2, p_2, c_2) \wedge p_1 \neq p_2)$$

$$\forall u, n_1, p_1, c_1, n_2, p_2, c_2 : \neg(\text{user}(u, n_1, p_1, c_1) \wedge \text{user}(u, n_2, p_2, c_2) \wedge c_1 \neq c_2)$$

## Part 1.3 Query Containment And Equivalence (Total: 36 Points)

### Question 1.3.1 (36 Points)

Consider the queries shown below. Check all possible containment relationships. If there exists a containment mapping from  $Q_i$  to  $Q_j$  then write down the mapping.

$Q_1(X, Y) :- R(X, Z), S(Z, A), R(A, Y).$

$Q_2(Y, X) :- R(Z, X), R(Y, B), R(C, B), R(D, B).$

$Q_3(X, Y) :- R(X, Z), R(A, Y).$

$Q_4(A, B) :- T(X, Z), R(A, X), R(Y, B), U(Z, Y).$

### Solution

$\frac{Q_1 \rightarrow Q_2:}{\text{no containment mapping exists}}$

$\frac{Q_1 \rightarrow Q_3:}{\text{no containment mapping exists}}$

$\frac{Q_1 \rightarrow Q_4:}{\text{no containment mapping exists}}$

$\frac{Q_2 \rightarrow Q_1:}{\phantom{\text{no containment mapping exists}}}$

$\frac{Q_2 \rightarrow Q_3:}{\phantom{\text{no containment mapping exists}}}$

$\frac{Q_2 \rightarrow Q_4:}{\phantom{\text{no containment mapping exists}}}$

$X \rightarrow Y$

$Y \rightarrow X$

$Z \rightarrow A$

$B \rightarrow Z$

$C \rightarrow X$

$D \rightarrow X$

$X \rightarrow Y$

$Y \rightarrow X$

$Z \rightarrow A$

$B \rightarrow Z$

$C \rightarrow X$

$D \rightarrow X$

$X \rightarrow B$

$Y \rightarrow A$

$Z \rightarrow Y$

$B \rightarrow X$

$C \rightarrow A$

$D \rightarrow A$

$\frac{Q_3 \rightarrow Q_1:}{\phantom{\text{no containment mapping exists}}}$

$\frac{Q_3 \rightarrow Q_2:}{\phantom{\text{no containment mapping exists}}}$

$\frac{Q_3 \rightarrow Q_4:}{\phantom{\text{no containment mapping exists}}}$

$X \rightarrow X$

$Y \rightarrow Y$

$Z \rightarrow Z$

$A \rightarrow A$

$X \rightarrow Y$

$Y \rightarrow X$

$Z \rightarrow B$

$A \rightarrow Z$

$X \rightarrow A$

$Y \rightarrow B$

$Z \rightarrow X$

$A \rightarrow Y$

$\frac{Q_4 \rightarrow Q_1:}{\text{no containment mapping exists}}$

$\frac{Q_4 \rightarrow Q_2:}{\text{no containment mapping exists}}$

$\frac{Q_4 \rightarrow Q_3:}{\text{no containment mapping exists}}$